

Why Types Matter

IN ALMOST 5 REASONS

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Silk 

<http://silkapp.com>

Silk 

One product

Started in 2009

Grown from 4 to 10 people

Located in Amsterdam

Why am I here?

We love functional programming!

Server: **Haskell**

Client: **JavaScript**



<http://haskell.org>

By default

Higher order

Lazy evaluation

Pure (no side effects)

Statically typed

Well suited for

Strict

Impure

Dynamically typed

Imperative

Hacking

5 REASONS

1.

CORRECTNESS

2.

REFACTORING

Types simplify refactoring

Particularly interesting:

Refactor into something
more type safe.

Employment

(Year, Year)

```
myProgram =  
  do (f, t) <- getEmployment "sebas"  
    printEmployment (t - f)
```

```
$ ./myProgram
```

```
Employment duration: -4 years.
```

Change representation?

(Year, Integer)

Change representation?

(Year, Unsigned)

Need to check every use site manually.

Create a new type

```
data Employment = MkEm Year Year
```



```
normalize f t = MkEm (min f t) (max f t)
```

```
make :: Year -> Year -> Maybe Employment  
make f t = if f < 2009  
           then Nothing  
           else Just (normalize f t)
```

```
module Employment  
  (Employment, make, from, to)
```

```
from, to :: Employment -> Year
```

```
from (MkEm f _) = f
```

```
to   (MkEm _ t) = t
```

Opaque datatype

Only export smart constructor `make`,

Not the original constructor `MkEm`.

Type error:

Couldn't match expected type **(Year, Year)**
with actual type **Employment**

```
myProgram =  
  do (f, t) <- getEmployment "sebas"  
    printEmployment (t - f)
```

```
myProgram =  
  do e <- getEmployment "sebas"  
     printEmployment (to e - from e)
```

```
$ ./myProgram
```

```
Employment duration: 4 years.
```

We win

Fixed a bug,
that will never occur again.

Made only a local change,
compiler points out use sites.

3.

REASONING

What does this function do?

```
foo :: [Bool] -> [Bool]
```

```
foo :: [Bool] -> [Bool]
```

The function can produce every
single bit sequence.

foo :: [a] -> [a]

foo :: [a] -> [a]

The function must reuse input.

reverse, empty, cycle, powerset, etc.

foo :: ()

Technically, not a function.

```
foo :: ()  
foo = ()
```

Singleton type 'unit',
only one value `()`.

foo :: a -> a


```
foo :: a -> a  
foo a = a
```

Only one possible implementation.

id

foo :: (a -> b) -> a -> b

```
foo :: (a -> b) -> a -> b
foo f a = f a
```

Function application, (\$)

(specialization of id)

foo ::= (b -> c)
-> (a -> b)
-> a -> c

foo :: (b -> c)
-> (a -> b)
-> a -> c
foo f g a = f (g a)

Function composition, $(.)$

foo :: (a, b) -> (b, a)

```
foo :: (a, b) -> (b, a)
foo (a, b) = (b, a)
```

Only one possible implementation.

swap

reverse :: [a] -> [b]

~~reverse :: [a] -> [b]~~

This is a lie!

Can only produce the empty list.

foo :: a

Maybe this?

foo :: a

foo = foo

Maybe this?

undefined :: a

Also called bottom, or \perp

foo :: a -> b

```
foo :: a -> b
```

Dangerous coercion!

Provided by the compiler as `unsafeCoerce`.

Theorems for free

Curry-Howard isomorphism:

Types \Leftrightarrow Propositions

Implementation \Leftrightarrow Proofs

4.

Genericity

Equality

```
equalInt :: Int -> Int -> Bool
```

Equality

```
equalInt :: Int -> Int -> Bool
```

```
equalStr :: String -> String -> Bool
```

```
equalBool :: Bool -> Bool -> Bool
```

```
equalX :: X -> X -> Bool
```

Generic Equality

```
(==) :: a -> a -> Bool
```

Nope

```
(==) :: a -> a -> Bool
```

Free theorems say cannot be done!

Do all types have equality anyway?

Constraint

```
(==) :: Eq a => a -> a -> Bool
```

Type classes

```
class Eq a where  
  (==) :: a -> a -> Bool
```

Type classes

```
instance Eq Bool where  
  True  == True  = True  
  False == False = True  
  _     == _     = False
```

Composability

```
notEq :: Eq a => a -> a -> Bool  
notEq a b = not (a == b)
```


Composability

```
lseq :: Eq a => [a] -> [a] -> Bool
```

```
lseq [] [] = True
```

```
lseq (x:xs) (y:ys) = x == y && lseq xs ys
```

```
lseq _ _ = False
```

Super classes

```
instance Eq a => Eq [a] where  
  (==) = \sEq
```

```
ghci> [[], [True, False]] == [[True]]
```

```
False
```

Type classes

Eq, Ord, Show, Read,
Random, Bounded, Enum,
IsString, Functor,
Num, Floating, Fractional,
Json, Xml, Binary, ...

Lots more

Deriving

```
data User = User
  { name      :: String
  , contact  :: Either Twitter Email
  , age      :: Int
  }
deriving (Eq, Ord, Show)
```

Either, Sum or +

```
data Either a b = Left a | Right b
```

Tuple, Product or *

```
data Tuple a b = Tuple a b
```

Tuple, Product or *

```
data (,) a b = (,) a b
```

```
data (a, b) = (a, b)
```


Deriving Generics

```
data User = User
  { name      :: String
  , contact  :: Either Twitter Email
  , age      :: Int
  }
deriving (Eq, Ord, Show, Generics)
```

Deriving Generics

Represents datatype algebraically,
using sums and products.

Now we can derive `Json`, `Xml`, `Binary`, etc.

5.

Effects

The Real World

```
getLine :: String
```

```
putStr :: String -> ()
```

```
readFile :: FilePath -> String
```

Put values in context

```
getLine :: IO String
```

```
putStr :: String -> IO ()
```

```
readFile :: FilePath -> IO String
```

Values

⋮ ⋮ a

becomes

⋮ ⋮ m a

Functions

$::: a \rightarrow b$

becomes

$::: a \rightarrow m b$

Effects

IO, ST, Cont,
Identity, Maybe, Either, [],
State, Reader, Writer,
Random, Parser,
Async, Par, STM, ...

Lots more

Running

```
runState :: State v a -> v -> (a, v)
```

Running

Most effects can be run

```
runState :: State v a -> v -> (a, v)
```

```
runReader :: Reader v a -> v -> (a)
```

```
runWriter :: Writer v a -> (a, v)
```

What about IO?

```
runIO :: IO a -> a
```

unsafePerformIO

~~runIO :: IO a -> a~~

Effects escape into purity!

```
main :: IO ()
```

The RTS interprets your top level `IO`.

Composability

Parse to
a list of values
or failures
from a socket
in parallel?

```
ParT . ParserT . ListT . EitherT Err . IO
```

Idioms

Category, Arrow,
Functor, Applicative,
Alternative,
Monad, MonadPlus,
Foldable, Traversable, ...

```
class Functor f where
```

```
  fmap :: (a -> b) -> f a -> f b
```



```
lengths :: Tree String -> Tree Int  
lengths = fmap length
```

class Applicative f **where**

pure :: a -> f a

(<*>) :: f (a -> b) -> f a -> f b

```
mkUser :: String -> Int -> User
```

```
pStr :: Parser String
```

```
pInt :: Parser Int
```

```
pUser :: Parser User
```

```
pUser = pure mkUser <*> pStr <*> pInt
```

```
class Foldable f where
```

```
  fold :: (a -> a -> a) -> f a -> a
```

```
minTree :: Tree Int -> Int  
minTree = fold min
```

```
class Monad m where
```

```
  return :: a -> m a
```

```
  (>>=)  :: m a -> (a -> m b) -> m b
```

```
getLine :: IO String
print   :: String -> IO ()
```

```
echo :: IO ()
echo = getLine >>= print
```

```
echo :: IO ()  
echo = do ln <- getLine  
         print ln
```



```
class Functor f => Traversable f where
```

```
mapM :: Monad m =>
```

```
(a -> m b) -> f a -> m (f b)
```

```
fetch :: [Request] -> ParT IO [Response]
fetch = mapM Http.request
```

```
fetch :: [Request] -> ParT IO [Response]
fetch = mapM Http.request
```

```
type Matrix a = Vector (Vector a)
```

```
transpose :: Matrix a -> Matrix a
```

```
transpose = mapM id
```

$$x:\sigma \in \Gamma$$

$$\Gamma \vdash x:\sigma$$

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